# A FAST INVERSE POLYNOMIAL RECONSTRUCTION METHOD BASED ON CONFORMAL FOURIER TRANSFORMATION

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**Abstract**—A fast Inverse Polynomial Reconstruction Method (IPRM) is proposed to efficiently eliminate the Gibbs phenomenon in Fourier reconstruction of discontinuous functions. The framework of the fast IPRM is modified by reconstructing the function in discretized elements, then the Conformal Fourier Transform (CFT) and the Chirp Z-Transform (CZT) algorithms are applied to accelerate the evaluation of reconstruction coefficients. The memory cost of the fast IPRM is also significantly reduced, owing to the transformation matrix being discretized in the modified framework. The computation complexity and memory cost of the fast IPRM are  $O(MN \log 2L)$  and O(MN), respectively, where L is the number of the discretized elements, M is the degree of polynomials for the reconstruction of each element, and N is the number of the Fourier series. Numerical results demonstrate that the fast IPRM method not only inherits the robustness of the Generalized IPRM (G-IPRM) method, but also significantly reduces the computation time and the memory cost. Therefore, the fast IPRM method is useful for the pseudospectral time domain methods and for the volume integral equation of the discontinuous material distributions.

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#### 1. INTRODUCTION

Traditional representation of a smooth and periodic function in the Fourier basis has been widely applied in the spectral method [1], the pseudospectral method [2, 3], and other computational physics [4] and signal processing research domains [5]. But Gibbs phenomenon occurs when the Fourier representation method is applied to a non-periodic or discontinuous function [6]. Neither accuracy nor efficiency can be achieved anymore because the finite Fourier representation cannot converge at the discontinuities or the boundaries of the function. Consequently, Gibbs phenomenon greatly hinders Fourier-representation based applications. A class of re-expansion methods have been developed to eliminate the Gibbs phenomenon in reconstructing discontinuous functions, with the locations of the discontinuities known a priori. They re-express the function in a nonperiodic basis set, such as the Gegenbauer polynomial basis set [6, 7], the Pade Fourier rational basis function [8] and the Freund polynomial basis set [9], to obtain the exponential convergence over the entire interval of the function, including the discontinuities or boundaries. However, the performance of these methods depends on the specific basis set and imposes several stringent requirements on the basis parameters [10, 11].

The Inverse Polynomial Reconstruction Method (IPRM) presents a basis-independent reconstruction frame to overcome the Gibbs phenomenon [12, 15, 17]. The IPRM reconstructs a function as a finite sum of polynomials, based on the equation between the Fourier series of the reconstructed function and that of the original function. The IPRM provides more flexibility and faster convergence than other reexpansion methods, but Jung and Shizgal found that the increment in the polynomial degree can dramatically deteriorate the conditioning of the transformation matrix and the reconstruction may fail to converge [13, 14].

The Generalized IPRM method (G-IPRM) has been proposed in [16] to avoid the ill-conditioning of the transformation matrix. The condition number of the transformation matrix in the G-IPRM is kept close to one by increasing the number of entries in the transformation matrix, specifically by increasing the number of Fourier series to be the square of the degree of polynomials. Thus, the well-conditioned transformation matrix makes the G-IPRM more robust, and significantly generalizes the application domains of the traditional IPRM, but with increased memory cost and computation complexity. The time complexity and the memory cost of the G-IPRM are both  $O(M_G^3)$ , where  $M_G$  is the degree of polynomials, if the iterative Least

Square QR (LSQR) method [18] is used to invert for the reconstruction coefficients. Such heavy computation complexity and memory cost are undesirable in many applications, e.g. the reconstruction of large-scale induced electric or current density in discontinuous media in computational electromagnetics [19–24] or in radar signal research domain [25, 26]. Therefore, it is necessary to reduce the computation time and the memory complexity in the G-IPRM.

In this paper, a fast IPRM method is presented to reduce the computation complexity and the memory cost of the G-IPRM. The framework of the fast IPRM is modified by reconstructing the function in the discretized elements, so that the Conformal Fourier Transform (CFT) [28] and the Chirp Z-Transform (CZT) [29,30] algorithms can be combined with the LSQR to reduce the computation complexity. The memory cost of the fast IPRM is reduced due to the transformation matrix being discretized, while the robustness of the reconstruction is preserved through configuring the relationships among the number of the discretized elements L, the degree of polynomials M and the number of Fourier series N. The computation complexity and memory cost of the fast IPRM are  $O(MN\log 2L)$  and O(MN), respectively. Numerical results demonstrate the robustness and the efficiency of the fast IPRM.

This paper is organized as follows: The fast IPRM method is presented in Section 2. Numerical results, including reconstruction of a meromorphic function and the induced current density in discontinuous material distributions, are shown in Section 3. The conclusions are drawn in Section 4.

#### 2. THE FAST IPRM METHOD.

Given the locations of discontinuities, the essence of the IPRM is solving an inverse problem from the already-known finite Fourier series of a discontinuous function and that of the polynomial basis functions. The objective of the proposed fast IPRM is to solve this inverse problem more efficiently through a discretization technique and two fast algorithms. There are three major steps in the fast IPRM: (1) modify the IPRM reconstruction framework by reconstructing the function in the discretized elements; (2) accelerate the reconstruction by combining two fast algorithms-CFT and CZT with the classic LSQR in the modified IPRM framework; (3) configure the discretization and reconstruction parameters to guarantee the robustness and efficiency enhancement of the fast IPRM.

In this section, first the fast IPRM will be presented for reconstructing a single-region discontinuous function;, and then it will

be extended to a multiple-region discontinuous function.

#### 2.1. The Fast IPRM Method for a Single-region **Discontinuous Function**

As a discontinuous function can be divided into multiple regions of piecewise continuous functions, we first consider the fast IPRM method for only one region of such a function.

#### 2.1.1. Modification of the IPRM Framework by Discretization

Assume that f(x) is continuous in the region  $\Omega = [a, b], b > a$  and zero outside.  $\mathbf{F} = [F_{-N/2}, \dots, F_{N/2-1}]$  is the vector of the first N Fourier coefficients of f(x) and the superscript T denotes matrix transpose.

The region  $\Omega$  is equally segmented into L uniformly-spaced elements  $\Omega = \bigcup_{l=1}^{L} \Omega_l$ , where  $\Omega_l = [a_l, b_l]$ , and  $\Omega_i \cap \Omega_j = \emptyset$  if  $i \neq j$ . The boundaries of the lth element are

$$a_l = a + (l-1)q, b_l = a + lq$$
 (1)

where q is the length of each element,  $q = b_l - a_l = \frac{b-a}{L}$ . The function f(x) is locally reconstructed in these uniformlydiscretized elements

$$\hat{f}_l(x) = \sum_{m=0}^{M-1} g_m^l \cdot C_m\left(\xi^l(x)\right), \qquad \hat{f}(x) = \bigoplus_{l=1}^L \hat{f}_l(x) \tag{2}$$

where  $\hat{f}_l(x)$  is the locally-reconstructed function in the lth element  $\Omega_l$ .  $\hat{f}(x)$  is the reconstructed function in the region  $\Omega$ , and  $\oplus$  denotes that  $\hat{f}(x)$  is a collection of  $\hat{f}_l(x)$ .  $\xi^l(x) = \frac{2x - (b_l + a_l)}{q}$  is the linear mapping from  $\Omega_l = [a_l, b_l]$  to [-1, 1]. M is the degree of the polynomial,  $C_m(\cdot)$  is the mth order polynomial, and  $\mathbf{g}^l = \begin{bmatrix} g_0^l & g_1^l & \dots & g_{M-1}^l \end{bmatrix}$  is the vector of coefficients for reconstructing  $f_l(x)$ .

The reconstruction coefficients can be obtained by equating the Fourier series of the reconstructed function  $\hat{f}(x)$  and that of the original function f(x)

$$\mathbf{W} \cdot \mathbf{g} = \mathbf{F} \tag{3}$$

where g is the vector of reconstruction coefficients constituted with that of the local reconstructions  $\mathbf{g} = \begin{bmatrix} \mathbf{g}^1 & \mathbf{g}^2 & \dots & \mathbf{g}^L \end{bmatrix}^T$ , and  $\mathbf{W}$  is the transformation matrix

$$\mathbf{W} = \left[ \begin{array}{cccc} \mathbf{w}^1 & \dots & \mathbf{w}^l & \dots & \mathbf{w}^L \end{array} \right] \tag{4}$$

where  $\mathbf{w}^l$  is the sub-matrix of Fourier series of polynomials for the local reconstruction  $\hat{f}(x)$ .

The entry of the sub-matrix  $\mathbf{w}^l$  at the n+N/2+1th row and the m+1th column is

$$w_{n,m}^{l} = \int_{a_{l}}^{b_{l}} C_{m} \left( \xi^{l} (x) \right) \cdot \exp \left( -i2\pi f_{0} nx \right) dx$$
$$= \exp \left[ -i2\pi n \left( l-1 \right) / L \right] \cdot w_{n,m}^{1}$$
(5)

where  $w_{n,m}^1 = \frac{q}{2} \cdot \exp\left[-i\pi n \left(a_1 + b_1\right) f_0\right] \cdot \int_{-1}^{1} C_m\left(x\right) \cdot \exp\left(-i\pi n x/L\right) dx$ , and  $f_0 = 1/(b-a)$  is the frequency resolution.  $n = -N/2, \ldots, N/2-1, m = 0, 1, \ldots, M-1, l = 1, \ldots, L$ , and i is the imaginary unit.

The modified IPRM framework for reconstructing single-region discontinuous function is achieved by Equations (1)–(5). The reconstruction framework of the G-IPRM is a special case with L=1. From (4) and (5), it can be seen that, due to the uniform discretization and the local reconstruction, the transformation matrix  $\mathbf{W}$  is uniformly discretized into L sub-matrices, which are differentiated only by the exponential scaling term independent of the polynomial order m. Given the Fourier series  $\mathbf{F}$ , the aim of the fast IPRM is to efficiently solve the reconstruction coefficients  $\mathbf{g}$  in (3) by taking advantage of the uniformly-discretized transformation matrix  $\mathbf{W}$  in the modified IPRM framework.

#### 2.1.2. Acceleration of the Reconstruction by the CFT and the CZT

When the iterative LSQR is utilized to solve the Equation (3), the computation complexity of the IPRM is determined by the number of the multiplication operations to calculate the following two matrix-vector products

$$\mathbf{y} = \mathbf{W} \cdot \mathbf{u} \tag{6}$$

$$\mathbf{z} = \mathbf{W}^{\dagger} \cdot \mathbf{v} \tag{7}$$

where the superscript † denotes the complex conjugate transpose,  $\mathbf{u}$  and  $\mathbf{v}$  are already known vectors,  $\mathbf{u} = [\mathbf{u}^1 \ \dots \ \mathbf{u}^l \ \dots \ \mathbf{u}^L]^T$ ,  $\mathbf{u}^l = [u^l_0 \ u^l_1 \ \dots \ u^l_{M-1}], \ \mathbf{v} = [v_{-N/2} \ \dots \ v_{N/2-1}]; \ \mathbf{y} \ \text{and} \ \mathbf{z} \ \text{are respectively} \ N \times 1 \ \text{and} \ LM \times 1 \ \text{dimension unknown vectors}.$ 

The computation complexities of (6) and (7) by direct multiplication are both O(LMN), which is very expensive. Fortunately, in the modified framework, taking advantage of the uniformly discretized transformation matrix  $\mathbf{W}$ , the computation complexity can be reduced by employing fast algorithms to evaluate products in (6) and (7).

Substituting (5) into (6), the (n + N/2 + 1)th entry of the vector **y** is

$$y_{n} = \sum_{l=1}^{L} \sum_{m=0}^{M-1} u_{m}^{l} \cdot w_{n,m}^{l} = \sum_{l=1}^{L} \int_{a_{l}}^{b_{l}} \gamma^{l}(x) \cdot \exp(-i2\pi n f_{0}x) dx$$
$$= \int_{a}^{b} \gamma(x) \cdot \exp(-i2\pi n f_{0}x) dx$$
(8)

where  $n = -N/2, \ldots, N/2 - 1$ ,  $\gamma^{l}(x)$  is the (M-1) order polynomial function in the lth element given by  $\gamma^{l}(x) = \sum_{m=0}^{M-1} u_{m}^{l} \cdot C_{m}(\xi^{l}(x))$ ,  $x \in [a_{l}, b_{l}].$ 

The product vector  $\mathbf{y}$  in (8) is the summation of Fourier series of the L piecewise continuous functions  $\gamma^1(x), \gamma^2(x), \dots, \gamma^L(x)$ . These functions are in the same polynomial degree M and have the same interval space q, thus they can be considered as the uniformly-spaced elements of a discontinuous polynomial function  $\gamma(x), \gamma(x) = \bigoplus_{l=1}^{L} \gamma^l(x)$  in [a, b] and zero outside.

Instead of calculating the Fourier series of the L piecewise continuous functions  $\gamma^1(x), \gamma^2(x), \ldots, \gamma^L(x)$  independently, the product vector  $\mathbf{y}$  can be efficiently and accurately achieved by calculating the Fourier series of  $\gamma(x)$  using the CFT algorithm [28], with the computation complexity of  $O(MN\log_2 2L)$ . The CFT algorithm, in which uniform discretization and high-order polynomial interpolation techniques are utilized, is a fast and accurate method to compute the Fourier series of a discontinuous function. Since the discretization is conformal to the discontinuities, the CFT algorithm can accelerate the evaluation of  $\mathbf{y}$ . Details of the CFT algorithm can be found in [28].

The other product vector to be calculated  $\mathbf{z}$  in (7) is composed of L sub-vectors  $\mathbf{z} = \begin{bmatrix} \mathbf{z}^1 & \dots & \mathbf{z}^l & \mathbf{z}^1 \end{bmatrix}^T$  and  $\mathbf{z}^l = \begin{bmatrix} z_0^l & z_1^l & \dots & z_{M-1}^l \end{bmatrix}$ ,  $l = 1, 2, \dots, L$ . Substituting (5) into (7), the (m+1)th entry of  $\mathbf{z}^l$  is

$$z_m^l = \sum_{n=-N/2}^{N/2-1} v_n \cdot \left(w_{n,m}^l\right)^* = \sum_{n=-N/2}^{N/2-1} d_m(n) \cdot \exp\left[i2\pi n \left(l-1\right)/L\right]$$
(9)

where  $d_m(n) = v_n \cdot \left(w_{n,m}^1\right)^*$  and the superscript \* denotes the complex conjugate.

Due to the transformation matrix **W** being uniformly discretized, for each constant integer  $m, z_m^1, z_m^2, \ldots, z_m^L$  are the first L discrete

Fourier series of the N-point discrete function  $d_m(\cdot)$ , with the frequency resolution  $f_0q = 1/L$  and  $L \ll N$ . Therefore, the product vector  $\mathbf{z}$  can be efficiently evaluated by performing the Chirp Z-Transform (CZT) [29,30] on the discrete function  $d_m(n)$ , for each m,  $m = 0, 1, \ldots, M - 1$ , hence the computation complexity to evaluate  $\mathbf{z}$  is  $O(MN \log_2 2L)$ .

By combining with the CZT and the CFT algorithms, the computation time in one iteration of the LSQR for evaluating the coefficient vector  $\mathbf{g}$  is reduced from O(LMN) to  $O(MN\log_2 2L)$ . Furthermore, it is not necessary to store the entire transformation matrix  $\mathbf{W}$ , and only the sub-matrix  $\mathbf{w}^1$  and several vectors are required, leading to a memory requirement of O(MN) in the fast IPRM algorithm.

#### 2.1.3. Discretization and Reconstruction Parameters

To guarantee the robustness as well as the efficiency enhancement of the fast IPRM, the relationship among the discretization and reconstruction parameters L, M and N are discussed.

The fast IPRM can possess the same robustness as the G-IPRM by keeping the condition number of the transformation matrix close to one. From (5), it can be seen that, the different sub-matrices of the transformation matrix  $\mathbf{W}$  are independent of each other. Therefore, the condition number of transformation matrix  $\mathbf{W}$  is equivalent to that of the matrix  $\mathbf{A}$ 

$$\mathbf{A} = \begin{bmatrix} H_0(-\pi N/2L) & \dots & H_{M-1}(-\pi N/2L) \\ \dots & \dots & \dots \\ H_0(\pi (N/2-1)/L) & \dots & H_{M-1}(\pi (N/2-1)/L) \end{bmatrix}$$
(10)

where  $H_m(\pi n/L) = \int_{-1}^1 C_m(x) \cdot \exp(-i\pi nx/L) dx$ . When  $C_m(\cdot)$  is the normalized Legendre polynomial basis set, the entry in matrix  $\mathbf{A}$  is  $H_m(\pi n/L) = (-i)^m (m+1/2)^{1/2} j_m(\pi n/L)$ , where  $j_m(\cdot)$  is the mth order spherical Bessel function of the first kind [27]. According to Theorem 4.1 and Theorem 4.2 in [16], the condition number of the matrix  $\mathbf{A}$  is decided by N/L. If N/L is the square of the degree M, i.e., if

$$N = LM^2 \tag{11}$$

then the condition number of the transformation matrix **W** is kept close to one. This not only preserves the robustness of the fast IPRM as that of the G-IPRM, but also ensures that the numbers of iterations of the two methods to evaluate **g** are close to each other once the termination criterion of the LSQR  $\varepsilon$  is determined [16]. Thus, the

reduction of the computation complexity in the fast IPRM is decided by that in one iteration of the LSQR.

To guarantee the efficiency improvement of the fast IPRM over the G-IPRM, the relationship among the reconstruction parameters of the fast IPRM L, M and that of the G-IPRM  $M_G$ , is now considered. As known from the convergence theory for polynomial expansion [1] and Theorem 5.1 in [16], the IPRM converges faster in a smaller interval. Because of the discretization and local reconstruction of the fast IPRM, given a certain accuracy requirement,  $M < M_G$ . Generally, a larger element number L accompanies with a lower polynomial degree M. For a given reconstruction accuracy requirement, L and M can be selected to satisfy

$$L^{2/3}M < M_G \tag{12}$$

From (12), the conditions  $LM^3 < L^2M^3 < M_G^3$  and  $LM^3\log_2 2L < L^2M^3 < M_G^3$  are satisfied. Therefore, with the same robustness and the same reconstruction accuracy requirement, the fast IPRM, whose computation complexity and memory cost are  $O\left(LM^3\log_2 L\right)$  and  $O\left(LM^3\right)$ , respectively, can achieve much higher computation and memory cost efficiency than the G-IPRM, whose oneiteration computation complexity and memory cost are both  $O\left(M_G^3\right)$ .

### 2.2. The Fast IPRM method for a Multiple-region Discontinuous Function

In this subsection, we will extend the fast IPRM to reconstruct a multiple-region discontinuous function from its Fourier coefficients.

Assume that f(x) is a multiple-region discontinuous function in the interval  $\Omega = [d_0, d_S]$ , with S > 1 and zero outside. The discontinuities of f(x) locate at  $[d_0, d_1, \ldots, d_S]$ , and  $d_0 < d_1 < \ldots < d_S$ . The first N Fourier series of f(x) is  $\mathbf{F} = \begin{bmatrix} F_{-N/2} \ldots F_{N/2-1} \end{bmatrix}^T$ .

To obtain the modified IPRM framework for reconstructing f(x), firstly the interval of the function  $\Omega$  is divided into S regions within which the function is continuous. Then by the same procedures as described in Subsection 2.1, each region is segmented into  $L_1, L_2, \ldots, L_S$  uniformly elements, where the function is locally reconstructed as the polynomials with degree of  $M_1, M_2, \ldots, M_S$ , respectively.

The transformation matrix **W**, with the size of  $N \times \sum_{s=1}^{S} L_s M_s$ , is

composed of  $\sum_{s=1}^{S} L_s$  sub-matrices of Fourier series of polynomial basis

for local reconstruction in each element

$$\mathbf{W} = \left[ \begin{array}{cccc} \mathbf{w}_1^1 & \dots & \mathbf{w}_s^l & \dots & \mathbf{w}_S^{L_S} \end{array} \right] \tag{13}$$

where  $\mathbf{w}_s^l$  is the submatrix of Fourier series, with the size of  $N \times M_s$ , for local reconstruction of the lth element in the sth region,  $s=1,\ldots,S$ ,  $l=1,\ldots,L$ . The entry of the sub-matrix  $\mathbf{w}_s^l$  at the n+N/2+1th row and the m+1th column is

$$w_{s,n,m}^{l} = \exp\left[-i2\pi n (l-1)f_0 q_s\right] \cdot w_{s,n,m}^{1}$$
(14)

where  $w_{s,n,m}^1 = \frac{q_s}{2} \cdot \exp[-i\pi n (2d_{s-1} + q_s)f_0] \cdot \int_{-1}^1 C_m(x) \cdot \exp(-i\pi n f_0 q_s x) dx$ .  $q_s = (d_s - d_{s-1})/L_s$  is element length in the sth region, and  $f_0 = 1/(d_s - d_0)$  is the frequency resolution.

The reconstruction of the multiple-region discontinuous function is also an inverse problem to evaluate the coefficients  $\mathbf{g}$  based on Equation (3). As known from (13) and (14), the transformation matrix  $\mathbf{W}$  is uniformly discretized in each region. Therefore, similar to the single-region case, the CFT and the CZT can be combined with the LSQR to accelerate the evaluation of the reconstruction coefficients, and the

computation complexity is  $O\left(N\sum_{s=1}^{S}M_{s}\log_{2}2L_{s}\right)$ . Meanwhile, the

memory cost is  $O\left(N\sum_{s=1}^{S}M_{s}\right)$  since only sub-matrices  $\mathbf{w}_{1}^{1}, \mathbf{w}_{2}^{1}, \ldots, \mathbf{w}_{S}^{1}$  and several vectors are required to be stored.

For the reconstruction of the multiple-region discontinuous function, the robustness and the efficiency of the fast IPRM are further guaranteed by selecting the reconstruction parameters to meet the following conditions

$$N = \max_{s=1}^{S} \left\{ M_s^2 / q_s \right\} (d_S - d_0)$$
 (15)

$$N\sum_{s=1}^{S} L_s M_s < N_G \sum_{s=1}^{S} M_{G,s}$$
 (16)

where  $M_{G,s}$  is the polynomial degree for reconstructing the sth region by the G-IPRM, with the same reconstruction accuracy as the fast IPRM, and  $N_G$  is the number of Fourier series in the G-IPRM,  $N_G = \max_{s=1}^{S} \left\{ M_{G,s}^2 / (d_s - d_{s-1}) \right\} \cdot (d_S - d_0)$ .

#### 3. NUMERICAL RESULTS

In this section, three examples are simulated to demonstrate the performance of the fast IPRM. Example in Subsection 3.1 takes the

same meromorphic function as in [16] to compare the fast IPRM and the G-IPRM proposed in [16]. Subsection 3.2 compares the reconstruction results of a single-region induced current density by the fast IPRM and the G-IPRM, respectively. Subsection 3.3 shows the reconstruction results by the fast IPRM and the G-IPRM for an induced current density whose support domain includes multiple piecewise continuous regions.

To compare the efficiency of the fast IPRM and the G-IPRM, in these three examples, the memory costs of the two methods are estimated by the number of the entries in the transformation matrix that need to be stored, and their computation complexities are measured by the number of the complex multiplication operations to obtain the reconstruction coefficients **g**.

### 3.1. Reconstruction of a Single-region Meromorphic Function

The analytic expression of the function is

$$f(x) = \begin{cases} 1/(x - 0.3i), & x \in [-1, 1] \\ 0, & \text{otherwise} \end{cases}.$$

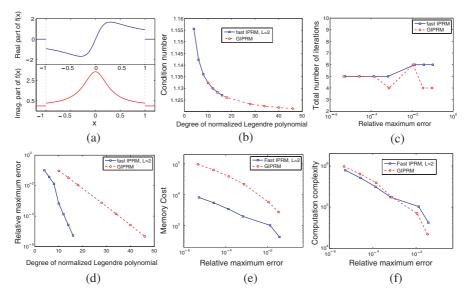
This function, as shown in Fig. 1(a), is a single-region discontinuous function with discontinuities located at  $x = \pm 1$ .

In the fast IPRM, for different degree of polynomials M, to meet condition (12), the number of the discretized elements is configured to be L=2. For the fixed L, according to (11), the number of the Fourier series is  $N=2M^2$ .

As shown in Fig. 1(b), similar to the G-IPRM, the condition number of the transformation matrix  $\mathbf{W}$  in the fast IPRM is kept close to one with the increasing of polynomial degree M. This not only ensures the robustness of the fast IPRM, but also makes the number of iterations in the fast IPRM almost the same as that in the G-IPRM when the termination criterion of the LSQR algorithm  $\varepsilon$  is decided, as shown in Fig. 1(c). Hence the difference of computation complexity between these two methods is determined by the computation complexity in one iteration.

The accuracy of the reconstruction is measured by the relative maximum error (RME)  $\left\|f(x) - \hat{f}(x)\right\|_{\infty} / \|f(x)\|_{\infty}$ . From Fig. 1(d), it can be seen that, the fast IPRM achieve the same reconstruction accuracy with a lower polynomial degree than the G-IPRM.

The memory cost and the computation complexity of the two methods are shown in Fig. 1(e) and Fig. 1(f), respectively. It can be seen that, although the computation complexity of the fast IPRM is not obviously improved, the memory cost of the fast IPRM is greatly



**Figure 1.** Reconstruction results of a Meromorphic function f(x) = 1/(x - 0.3i). (a) Function f(x). (b) Condition number of transformation matrix in the fast IPRM. (c) Total iteration number required in the G-IPRM and in the fast IPRM to meet the termination criterion  $\varepsilon = 10^{-13}$  in the LSQR. (d) Relative maximum error of reconstruction by the G-IPRM and by the fast IPRM. (e) The memory cost of the fast IPRM and the G-IPRM. (f) The computation complexity of the fast IPRM and the G-IPRM.

reduced, e.g., for a relative maximum error as  $10^{-5}$ , the memory cost in the fast IPRM is about 1/12 of that in the G-IPRM.

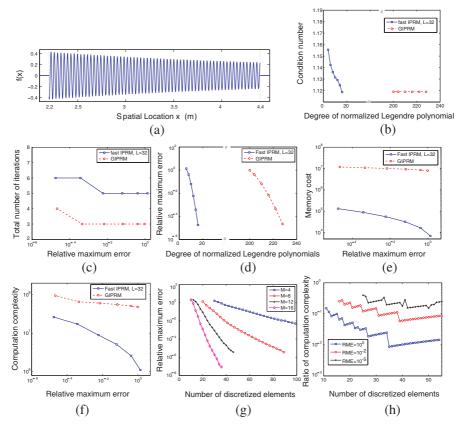
Therefore, for this single-region meromorphic function, the fast IPRM not only inherits the robustness of the G-IPRM but also possesses higher memory efficiency than the G-IPRM.

## **3.2.** Reconstruction of a Single-region Induced Current Density

In this subsection, an induced electric current density due to a plane wave at normal incidence to a three-layer medium with frequency  $f_c=5\,\mathrm{GHz}$  in a vacuum background is reconstructed. The layer interfaces are located at 2.2 m and 4.4 m, respectively, and the relative magnetic permeability for all three layers is 1. The relative permittivity is  $\epsilon_r=1$ , 3 and 1 respectively. The conductivity for the center layer is  $\sigma=0.003\,\mathrm{S/m}$ . The wave number in the central layer is 63.5085. As shown in Fig. 2(a), this wave distribution is a single-

region discontinuous function with discontinuities located at  $x=2.2\,\mathrm{m}$  and  $x=4.4\,\mathrm{m}$ , respectively.

To guarantee the efficiency improvement of the fast IPRM, L and M are selected to satisfy (12). As shown in Fig. 2(h) multiple values of L can meet formula (12) and reduce the computation complexity



**Figure 2.** Reconstruction results of an induced electric current density in a three-layer medium. (a) Spatial distribution of an induced electric current density in a three-layer medium. (b) Condition number of transformation matrix in the fast IPRM. (c) Total iteration number required in the G-IPRM and in the fast IPRM to meet the termination criterion  $\varepsilon=10^{-6}$  in the LSQR. (d) Relative maximum error of reconstruction by the G-IPRM and by the fast IPRM. (e) The memory cost of the fast IPRM and the G-IPRM. (f) The computation complexity of the fast IPRM and the G-IPRM. (g) The relative maximum error of the fast IPRM. (h) The ratio of computation complexity between the fast IPRM and the G-IPRM.

compared to the G-IPRM. In the simulation, the value of L is taken as the half of the wave numbers in the center layer, i.e., L=32, and for the fixed L, the number of the Fourier series varies with the degree of polynomial,  $N=32M^2$ .

Figure 2(b) shows that, just as in the G-IPRM, the condition number of the transformation matrix  $\mathbf{W}$  in the fast IPRM is kept close to one with the increasing of polynomial degree. Fig. 2(c) demonstrates that the numbers of iterations in the fast IPRM and the G-IPRM are close to each other, with a fixed termination criterion  $\varepsilon$  in the LSQR. Thus the difference of computation complexity between these two methods is determined by that in one iteration of the LSQR.

As shown in Fig. 2(d), the G-IPRM method converges much slower and requires higher degree polynomials, whereas the fast IPRM with elements number L=32, needs a much lower polynomial degree to obtain the same reconstruction accuracy. Fig. 2(g) shows that the RME of the fast IPRM is reduced with the increase of L for a fixed value of M. Meanwhile, a larger element number L accompanies a lower polynomial degree M for a given accuracy requirement.

The memory cost and the computation complexity are significantly reduced in the fast IPRM compared with the G-IPRM, as shown in Figs. 2(e) and (f). With the relative maximum error  $10^{-5}$ , the computation complexity and the memory cost in fast IPRM are only about 1/4 and 1/90 of that in the G-IPRM, respectively.

## **3.3.** Reconstruction of a Multiple-region Induced Current Density

A multiple-region induced electric current density due to a plane wave at normal incidence to a five-layer medium with frequency  $f_c = 5\,\mathrm{GHz}$  in a vacuum background, is reconstructed in this subsection. The layer interfaces are located at  $0.8\,\mathrm{m}$ ,  $1\,\mathrm{m}$ ,  $5.4\,\mathrm{m}$  and  $5.8\,\mathrm{m}$ , respectively. The relative permittivity is  $\epsilon_r = 1, 5, 3, 4$  and 1, respectively. The conductivity for the center three layers is  $\sigma = 0.005, 0.003$  and  $0.004\,\mathrm{S/m}$ , respectively, and the wave numbers in the three center layers are 7.4536, 127.02, 13.3333, respectively. As shown in Fig. 3(a), the wave distribution is a three-region discontinuous function with discontinuities located at  $0.8\,\mathrm{m}$ ,  $1\,\mathrm{m}$ ,  $5.4\,\mathrm{m}$  and  $5.8\,\mathrm{m}$ , respectively.

In the fast IPRM, to meet the condition in (16), we configure the element number as the half of the wave numbers in the central three layers,  $L_1 = 4$ ,  $L_2 = 64$  and  $L_3 = 7$ . The degrees of polynomials in the three regions are the same  $M_1 = M_2 = M_3 = M$ , and according to (15), the number of the Fourier series is  $N = \max_{s=1}^{S} \{1/q_s\} (d_S - d_0) M^2 = 100 M^2$ .

As in the previous two numerical examples, the robustness and the efficiency of the fast IPRM are illustrated in Fig. 3(b)–Fig. 3(f). Fig. 3(b) shows that, in the fast IPRM, the condition number of transformation matrix **W** is kept close to one with the increasing of polynomial degree, while that in the G-IPRM is slightly larger.

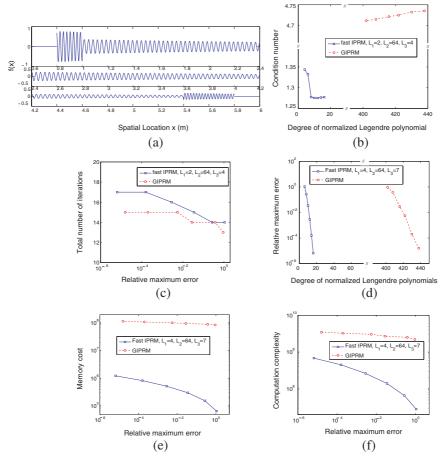


Figure 3. Reconstruction results of an induced electric current density in a five-layer medium. (a) Spatial distribution of an induced electric current density in a five-layer medium. (b) Condition number of transformation matrix in the fast IPRM. (c) Total iteration number required in the G-IPRM and in the fast IPRM to meet the termination criterion  $\varepsilon=10^{-13}$  in the LSQR. (d) Relative maximum error of reconstruction by the G-IPRM and by the fast IPRM. (e) The memory cost of the fast IPRM and the G-IPRM. (f) The computation complexity of the fast IPRM and the G-IPRM.

Fig. 3(c) demonstrates that the numbers of iterations in the fast IPRM and the G-IPRM are close to each other. As shown in Fig. 3(d), the polynomial degree in the fast IPRM is much lower than that in the G-IPRM. From Fig. 3(e) and Fig. 3(f), it can be found that, the memory cost and the computation complexity are significantly reduced in the fast IPRM compared with the G-IPRM, e.g. for a maximum relative error as  $10^{-5}$ , the computation complexity and the memory cost in fast IPRM are about 1/5 and 1/94 of that in the G-IPRM, respectively.

#### 3.4. Discussion

The above numerical examples demonstrate that the fast IPRM method works efficiently for eliminating the Gibbs phenomenon of both single-region and multiple-region discontinuous functions. It is also seen that, provided that the number of the discretized elements L satisfy (12) or (16), the efficiency improvement of the fast IPRM over the G-IPRM is more distinct with the increment of L. For the multiple-region case, the improvement of the efficiency in the fast IPRM over the G-IPRM is decided by the region, which has the largest number of the discretized elements.

#### 4. CONCLUSION

In this paper, we propose a fast IPRM method to efficiently overcome Gibbs phenomenon in the reconstruction of discontinuous or non-periodic functions. The Conformal Fourier Transform (CFT) and the Chirp Z-Transform (CZT) algorithms are combined in the modified IPRM framework to decrease both computation time and memory complexity. For a discontinuous function with S piecewise-continuous regions, the memory cost and the computation complexity of the

fast IPRM method are 
$$O\left(N\sum_{s=1}^{S}M_{s}\right)$$
 and  $O\left(N\sum_{s=1}^{S}M_{s}\log_{2}2L_{s}\right)$ ,

respectively. Numerical results show that the fast IPRM method works well for both single-region functions and multiple-region functions. Therefore, the fast IPRM method can be useful in pseudospectral time-domain methods and in volume integral equation solvers for discontinuous material distributions or other areas where Gibbs phenomenon needs to be efficiently eliminated.

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